## IN THE CLAIMS

Please amend the claims to be in the form as follows:

Claim 1 (original): A method of determining a best-case response time of a first periodic task, the method comprising:

a first step of determining that the first periodic task has a lower priority than a higher priority of a second periodic task,

characterized in that the method further comprises:

a second step of determining that the best-case response time of the first periodic task is substantially equal to the difference between a start of the first periodic task and a completion of the first periodic task, the start being right after a release of the first periodic task and the completion coinciding with a release of the second periodic task.

Claim 2 (original): A method of determining a best-case response time of a first periodic task according to claim 1, wherein  $BR_i$  denotes the best-case response time of the first periodic task,  $BR_i$  being substantially equal to the largest value that satisfies:

$$BC_{i} + \sum_{j \in h_{p}(i)} \left( \left\lceil \frac{BR_{i}}{T_{j}} \right\rceil - 1 \right) BC_{j}$$

wherein  $BC_i$  denotes a best-case computation time of the first periodic task  $\tau_i$ , hp(i) denotes a set of tasks with a higher priority than the lower priority,  $T_j$  denotes a period of activation of a task j of hp(i).

Claim 3 (original): A method of determining a best-case response time of a first periodic task according to claim 2, wherein  $WR_i$  denotes a worst-case response time of the first periodic task  $\tau_i$  and the best-case response time  $BR_i$  can be found by an iterative procedure of k iterations, where  $k = 0, 1, \dots$  comprising:

$$BR_i(0) = WR_i$$

$$BR_{i}(k+1) = BC_{i} + \sum_{j \in lop(i)} \left( \left\lceil \frac{BR_{i}(k)}{T_{j}} \right\rceil - 1 \right) BC_{j}$$

wherein the iterative procedure terminates when the same value is found for two successive iterations of k.

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Claim 4 (original): A method of determining a best-case response time of a first periodic task according to Claim 3, wherein the worst-case response time of the first periodic task is based upon a worst-case computation time of the first periodic task.

Claim 5 (original): A method of determining a best-case response time of a first periodic task according to Claim 3, wherein the worst-case response time of the first periodic task is based upon a best-case computation time of the first periodic task.

Claim 6 (original): A method of determining a best-case response time of a first periodic task according to claim 3, wherein  $RJ_i$  denotes a release jitter of the first periodic task  $\tau_i$ , the release jitter being a variation in the release of the first periodic task and the release jitter having a negative contribution to the best-case response time:

$$BR_i(0) = WR_i$$

$$BR_{i}(k+1) = BC_{i} + \sum_{i \in ho(i)} \left( \left\lceil \frac{BR_{i}(k) - RJ_{i}}{T_{i}} \right\rceil - 1 \right)^{+} BC_{i}$$

wherein  $x^+$  denotes the maximum of 0 and x.

Claim 7 (original): A system for determining a best-case response time of a first periodic task, the system comprising:

determination means conceived to determine that the first periodic task has a lower priority than a higher priority of a second periodic task, characterized in that the system further comprises:

response time means conceived to determine that the best-case response time of the first periodic task is substantially equal to the difference between a start of the first periodic task right after its release and a completion of the first periodic task that coincides with a release of the second periodic task.

Claim 8 (original): A system (800) of determining a best-case response time of a first periodic task according to claim 7, the system further comprising first calculation means (802) conceived to calculate the best-case response time denoted by  $BR_i$  according to the following formula:

$$BR_{i} = BC_{i} + \sum_{j \in ho(i)} \left( \left[ \frac{BR_{i}}{T_{j}} \right] - 1 \right) BC_{j}$$

wherein  $BC_i$  denotes a best-case computation time of the first periodic task  $\tau_i$ , hp(l) denotes a set of tasks with a higher priority than the priority of the first periodic task,  $T_j$  denotes a period of activation of a task j of hp(i), and  $BR_i$  denotes the best-case response time of the first periodic task.

Claim 9 (original): A system (800) of determining a best-case response time of a first periodic task according to claim 8, the system further comprising second calculation means (804) conceived to calculate the best-case response time denoted by  $BR_i$  according to the following iterative procedure of k iterations, where k = 0.1...:

$$BR_i(0) = WR_i$$

$$BR_i(k+1) = BC_i \cdot \sum_{f \in hp(i)} \left( \left\lceil \frac{BR_i(k)}{T_i} \right\rceil - 1 \right) BC_i$$

wherein  $WR_i$  denotes a worst-case response time of the first periodic task  $\tau_i$  and the iterative procedure terminates when the same value is found for two successive iterations of k.

Claim 10 (original): A system (800) of determining a best-case response time of a first periodic task according to claim 7, the system further comprising third calculation means (806) conceived to calculate the best-case response time denoted by  $BR_i$  corrected for a release jitter, the release jitter being a variation in the release of the first periodic task:

$$BR_i(0) = WR_i$$

$$BR_{i}(k+1) = BC_{i} + \sum_{j \in loc(i)} \left( \left\lceil \frac{BR_{i}(k) - RJ_{j}}{T_{j}} \right\rceil - 1 \right)^{+} BC_{j}$$

wherein  $RJ_i$  denotes the release jitter of the first periodic task  $\tau_i$ , and  $x^+$  denotes the maximum of 0 and x.

Claim 11 (previously presented): A computer program product arranged to perform the method according to Claim 1.

Claim 12 (original): A storage device (812) comprising a computer program product according

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to Claim 11.

Claim 13 (previously presented): A television set (910) comprising a system according to Claim 7.

Claim 14 (previously presented): A set-top box (1002) comprising a system according to Claim 7.

Claim 15 (new): The set-top box of Claim 14 wherein system determines from the best-case response time of the first periodic task and a worst case response time of the first periodic task if a signal can be used by the system.

Claim 16 (new): The television set of Claim 13 wherein system determines from the best-case response time of the first periodic task and a worst case response time of the first periodic task if a signal can be used by the system.

Claim 17 (new): A set-top box arranged to perform the method according to Claim 1 wherein the method determines from the best-case response time of the first periodic task and a worst case response time of the first periodic task if a signal can be used by the set-top box.

Claim 18 (new): The television set arranged to perform the method according to Claim 1 wherein method determines from the best-case response time of the first periodic task and a worst case response time of the first periodic task if a signal can be used by the television set.